A Machine-Independent Port of the MPD Language Run Time System to NetBSD Operating System

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1 Introduction

MPD (presented in Gregory Andrews' book about Foundations of Multithreaded, Parallel, and Distributed Programming[1]) is the successor of SR[2] ("synchronizing resources"), a PASCAL–style language enhanced with constructs for concurrent programming developed at the University of Arizona in the late 1980s[3].

MPD as implemented provides the same language primitives as SR with a different syntax which is closer to C.

The run-time system (in theory, identical) of both languages provides the illusion of a multiprocessor machine on a single single— or multi–CPU Unix–like system or a (local area) network of Unix-like machines.

Chair V of the Computer Science Department of the University of Bonn is operating a laboratory for a practical course in parallel programming consisting of computing nodes running NetBSD/arm, normally used via PVM, MPI, etc.

We are considering to offer SR and MPD for this, too. As the original language distributions are only targeted at a few commercial Unix systems, some porting effort is needed, outlined in the SR porting guide[7] and also applicable to MPD.

The integrated POSIX threads support of NetBSD-2.0 enables us to use library primitives provided for NetBSD's pthread system to implement the primitives needed by the SR and MPD run-time systems, thus implementing 13 target CPUs with a one-time effort; once implemented, symmetric multiprocessing (SMP) would automatically be used on any multiprocessor machine with VAX, Alpha, PowerPC, Sparc, 32-bit Intel and 64 bit AMD CPUs.

This paper describes mainly the MPD port. Porting SR was started earlier and partially described in [6] (Assembler and SVR4 cases) while only preliminary results for our new approach could be presented at the conference.

Most of the differences between our changes to SR and to MPD could be done by mechanically replacing mpd_ by sr_ in the code; because of this, and because the

test machine	A	В
architecture	i386	arm
CPU	Pentium 4	SA-110
clock	$1600~\mathrm{MHz}$	233 MHz
cache	2 MB	16kB I + 16 kB D

Table 1: Test machines

Implementation	A	В	
assembler	$0.013 \ \mu s$	n/a	
context_u libary calls	$0.138 \ \mu s$	$0.237 \; \mu { m s}$	
SVR4 system calls	$1.453 \; \mu { m s}$	$9.649 \ \mu s$	

Table 2: Raw context switch times

machine-independent parts of the SR and MPD run-time support are identical (according to the authors) all results (especially timing results) equally apply to the SR port. (This has been verified.)

2 Generic Porting Problems

Despite the age of SR, the latest version (2.3.3) had been changed to use <stdarg.h> instead of <varargs.h>, thus cutting the number of patches needed for NetBSD 2.0 and later by half compared to the original porting effort described in [6]. MPD 1.0.1 contains no traces of <varargs.h>.

The only patches – outside of implementing the context switching routines – were for 64 bit cleanliness (see also [5]).

3 Verification methods

MPD itself provides a verification suite for the whole system; also a small basic test for the context switching primitives. There is no split between the basic and the extended verification suite, as in SR.

3.1 Context Switch Primitives

The context switch primitives can be independently tested by running make in the subdirectory csw/ of the distribution; this builds and runs the cstest program, which implements a small multithreaded program and checks for detection of stack overflows, stack underflows, correct context switching etc.[7] This test is automatically run when building the whole system.

3.2 Overall System

When the context switch primitives seem to work individually, they need to be tested integrated into the run-time system. The SR and MPD authors provide a verification suite in the vsuite/ subdirectory of the distributions to achieve this, as well as testing the the building system used to build MPD, and the mpd compiler, mpd1 linker, etc.

It is run by calling the driver script mpdv/mpdv, which provides options for selecting normal vs. verbose output, as well as selecting the installed vs. the freshly compiled MPD system.

For all porting methods described below (assembler primitives, SVR4 system calls and NetBSD pthread library calls), the full verification suite has been run and any reported problem has been fixed.

Test description	i386 ASM	context_u	SVR4 s.c.
loop control overhead	$0.002~\mu {\rm s}$	$0.002 \; \mu { m s}$	$0.002 \; \mu s$
local call, optimised	$0.011 \ \mu s$	$0.011 \; \mu { m s}$	$0.011 \; \mu { m s}$
interresource call, no new process	$0.270 \; \mu { m s}$	$0.260 \; \mu { m s}$	$0.250 \; \mu { m s}$
interresource call, new process	$0.650 \; \mu { m s}$	$4.200 \; \mu { m s}$	$4.350 \; \mu { m s}$
process create/destroy	$0.540 \; \mu { m s}$	$4.020 \ \mu s$	$4.280 \; \mu { m s}$
semaphore P only	$0.011 \; \mu { m s}$	$0.011~\mu {\rm s}$	$0.011 \; \mu { m s}$
semaphore V only	$0.008 \; \mu { m s}$	$0.008 \; \mu { m s}$	$0.008 \; \mu { m s}$
semaphore pair	$0.019 \; \mu { m s}$	$0.019 \ \mu s$	$0.019 \ \mu s$
semaphore requiring context switch	$0.110 \ \mu s$	$0.220 \; \mu { m s}$	$1.550 \ \mu s$
asynchronous send/receive	$0.300 \; \mu s$	$0.290 \; \mu { m s}$	$0.300 \; \mu s$
message passing requiring context switch	$0.400 \; \mu { m s}$	$0.560 \; \mu { m s}$	$1.920 \ \mu s$
rendezvous	$0.600 \; \mu { m s}$	$0.850 \; \mu s$	$4.200 \ \mu s$

Table 3: Run time system performance, system A (Pentium 4, 1600 MHz). The median times reported by the MPD script vsuite/timings/report.sh are shown.

4 Performance evaluation

MPD comes with two performance evaluation packages. The first, for the context switching primitives, is in the csw/subdirectory of the source distribution; after make csloop you can start ./csloop N where N is the number of seconds the test will run approximately.

Tests of the language primitives used for multithreading are in the vsuite/timings/subdirectory of the source tree enhanced with the verification suite. They are run by three shell scripts used to compile them, executed them, and summarize the results in a table.

5 Establishing a baseline

There are two extremes possible when implementing the context switch primitives needed for MPD: implementing each CPU manually in assembler code (what the MPD implementation does normally) and using the SVR4-style functions getcontext(), setcontext() and swapcontext() which operate on struct ucontext; these are provided as experimental code in the file csw/svr4.c of the MPD distribution.

The first tests were done by using the provided i386 assembler context switch routines. After verifying correctness and noting the times (see tables 2 and 3), the same was done using the SVR4 module instead of the assembler module.

These tests were done on a Pentium 4 machine running at 1600 MHz with 2 megabytes of secondary cache, and 1 GB of main memory, running NetBSD-3.0_BETA as of end of October 2005.

The SVR4 tests were redone on a DNARD system (for its ARM cpu, no assembler stubs are provided in either the SR or MPD distributions).

Table 3 shows a factor-of-about-ten performance hit for the operations that require context switches; note, however, that the absolute values for all such operations are still smaller than $5\,\mu s$ on 1600 MHz machine and will likely not be noticeable if a parallelized program is run on a LAN-coupled cluster: on the switched LAN connected to the test machine, the time for an ICMP echo request to return is about 200 μs .

Test description	ARM ASM	context_u	SVR4 s.c.
loop control overhead	n/a	$0.057 \; \mu { m s}$	$0.056 \; \mu { m s}$
local call, optimised	n/a	$0.376 \; \mu { m s}$	$0.355 \; \mu { m s}$
interresource call, no new process	n/a	$4.300 \ \mu s$	$4.080 \; \mu { m s}$
interresource call, new process	n/a	$27.250 \ \mu s$	$55.900 \ \mu s$
process create/destroy	n/a	$25.240 \ \mu s$	$58.780 \ \mu s$
semaphore P only	n/a	$0.304~\mu {\rm s}$	$0.301 \; \mu s$
semaphore V only	n/a	$0.254 \ \mu s$	$0.249 \ \mu s$
semaphore pair	n/a	$0.506 \; \mu { m s}$	$0.487~\mu {\rm s}$
semaphore requiring context switch	n/a	$1.570 \ \mu s$	$11.180 \ \mu s$
asynchronous send/receive	n/a	$5.550 \ \mu s$	$5.190 \ \mu s$
message passing requiring context switch	n/a	$6.740 \ \mu s$	$30.140 \ \mu s$
rendezvous	n/a	$9.600 \; \mu s$	$54.000 \ \mu s$

Table 4: Run time system performance, system B (StrongARM SA-110, 233 MHz). The median times reported by the MPD script vsuite/timings/report.sh are shown.

6 Improvements using NetBSD library calls

While using the system calls getcontext and setcontext, as the svr4 module does, should not unduly penalize an application distributed across a LAN, it might be noticeable with local applications.

However, we should be able to do better than the svr4 module without writing our own assembler modules, since NetBSD 2.0 (and later) contains its own set of them for the benefit of its native Posix threads library (libpthread), which does lots of context switches within a kernel provided light weight process[8]. The primitives provided to libpthread by its machine dependent part are the three functions _getcontext_u, _setcontext_u and _swapcontext_u with similar signatures as the SVR4-style system calls getcontext, setcontext and swapcontext.

There were a few difficulties that arose while pursuing this.

First, on one architecture (i386) _setcontext_u and _getcontext_u are implemented by calling through a function pointer which is initialized depending on the FPU / CPU extension mode available on the particular CPU used (8087-mode vs. XMM). On this architecture, _setcontext_u and _getcontext_u are defined as macros in a private header file not installed. The developer in charge of the code has indicated that he might implement public wrappers; until then, we have to check all available NetBSD architectures and copy the relevant code to our module csw/netbsd.c.

Second, we need to extract the relevant object modules from the threading library for static linking (libpthread.a) without resolving any other symbols, because normal libpthread is overloading some system calls thus causing failure of applications not properly initializing it.

Again, this set of context switch code has been verified by running cstest and the full verification suite.

The low-level as well as the high-level timings with the new context switch package have again been collected in tables 2, 3 and 1.

To ease installation, a package for the NetBSD package system has been built for SR and MPD, available in the lang/sr and lang/mpd subdirectories of the pkgsrc root.

As the NetBSD package system is available for more operating systems than Net-BSD[4], a lot more work would be needed to make the packages universal; thus they are restricted to be built on NetBSD 2.0 and later.

7 Discussion

Our new approach has raw context switch times that are only 10% of the SVR4 system call ones. Compared to the assembler routines, they are only slower by a factor of 10 (see table 2).

Table 3 shows three classes of high level operations.

- 1. Non-context switching operations have the same speed independent of the context switch primitives used, as expected.
- 2. The two operations measured requiring a process creation (in the MPD language sense) are about as fast as in the SVR4-system-call case. This was expected, as the process creation primitive does a system call internally.
- 3. Context switching operations which do not create a new process (in the MPD language sense) are slower than in the assembler case, but faster than in the SVR4-style case, by an amount roughly equivalent to one (semaphore operation, message passing) or two (rendezvous) context switching primitive times.

The same classification can be done for the 233 MHz ARM CPU (table 1). However, SVR4 process creation, destruction and the rendezvous need about one third of the LAN two-way network latency, thus cannot be neglected anymore. We conclude that for machines in the 300 MHz range and below, using assembler implementation (where available) or at least our new implementation of the context switching primitives is a necessity. This is also expected for even slower machines.

MPD can be compiled in a mode where it will make use of multiple threads provided by the underlying OS, so that it can use more than one CPU of a single machine. This has not been implemented yet for NetBSD, but should be.

8 Summary

A method for porting SR and MPD to NetBSD has been shown, for which only preliminary results, and only for SR, were presented earlier.

The SR porting effort was easily adopted for the MPD case. In fact, the run time system (library and srx/mprx) could probably be factored out into a common run-time system package.

The new port was verified using the SR and MPD verification suites.

As discussed above, the SVR4-system-call approach, while feasible, creates an overhead that is clearly visible for non-networked operation of a distributed program; on our Pentium machine, high level context switching operations are slower by a factor between 7 and 11 (the raw context switch primitives are slower by a factor of 110). Even for networked operation, for a 233 MHz StrongArm CPU or slower machines, context switch latency exceeds one third of the network latency.

The approach using the libpthread primitives is much faster for all but the process creation/destruction case and should thus be adequate for about any application in the networked case, and for any in the single-machine case that does not do excessive amounts of implicit or explicit process creation.

For highly communication-bound problems on a single machine, using the assembler primitives might show a visible speedup, where available.

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